Lecture #5

云计算入门 Introduction to Cloud Computing GESC1001

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Course schedule

Part I	Introduction and overview
Part 2	Distributed and parallel systems
Part 3	Cloud infrastructure
Part 4	Cloud application paradigm (1)
Part 5	Cloud application paradigm (2)
Part 5 Part 6	Cloud application paradigm (2) Cloud virtualization and resource management
Part 5 Part 6 Part 7 & 8	Cloud application paradigm (2) Cloud virtualization and resource management Cloud computing storage systems Cloud computing security



Last week:

- Review
- **Chapter 4**: Cloud application paradigm (part I)

Today:

- Chapter 4: Cloud applications (part 2) – the Map Reduce model
- Assignment I



How to ask questions

We can discuss immediately after lectures

You may use the QQ group to contact teaching assistants

My e-mail: philfv8@yahoo.com

4-CLOUD APPLICATIONS (云应用) PART 2

- We discussed challenges for developing cloud applications
- Today, we will talk about the details of how cloud applications are created.
- To make **cloud applications**, the **MapReduce model** is very popular.
- It is a "**programming model" (**编程模型 a way of developing applications for the cloud).
- It was proposed by Google in a research paper, published in 2004.

MapReduce: Simplified Data Processing on Large Clusters, Jeffrey Dean and Sanjay Ghemawat, OSDI'04: Sixth Symposium on Operating System Design and Implementation, San Francisco, CA, December, 2004.

Why MapReduce is popular?

- Because it is a simple programming model.
- A programmer (程序员) can easily write an application that run on a distributed system (the cloud), without much experience about how distributed systems work.





- One of the most popular version of MapReduce is Hadoop.
- It is an open-source (开放源码) implementation of MapReduce.
- I will explain the main idea.
- We will also discuss **three examples**.





- The main advantage of the cloud is elasticity (云的弹性).
 - Using as many computers as needed to address the cost (元) and timing constraints of an application.
 - Sharing the workload (工作负载) between several computers.
 - It must be divided into sub-tasks that can be accomplished in parallel by several computers.
- But how to do this?



The **workload** should be divided (分配) approximately equally between computers.

Workload



- Partitioning (分配) the workload is not always easy.
- Three main types of workloads:
 - modularly divisible (模块化分割)
 workload: the workload is already divided into sub-tasks.
 - arbitrarily divisible (可任意划分) workload: the workload can be partitioned into an arbitrarily large number of sub-tasks of equal or similar size.
 - Others.

Map Reduce

- Designed for **arbitrarily divisible** (可任意划分) **workloads**.
- It is used to perform parallel processing (并行处理) for data-intensive (数据密集型) applications.
- It has many applications: e.g. physics, biology, etc.
- Once a cloud applications is created using MapReduce, it can run in the cloud on as many computers as needed.

Phase I (Map)

- I. Split the data into blocks
- 2. Assign each block to an instance (实例) (e.g. a computer or virtual machine)
- 3. Run these instances in parallel



Data

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Phase 2 (Reduce)

- 1. Once all the instances have finished their subtasks, they send their results.
- 2. Results are merged to obtain the final result.



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Phase 2 (Reduce)

Result

- I. Once all the instances have finished their subtasks, they send their results.
- 2. Results are merged to obtain the final result.



Who split the data and the gather results?

• A "master instance" takes care of splitting the data.



 Merging the results can be done by a set of instances called the "reducing instances"



How data is represented?

The **input data** (输入数据) can be any kind of files.

But it is converted to a set of <key, value> pairs (键值对).

e.g.: (key= CN, value = Shenzhen) (key= CN, value = Beijing)

A key (()) is some information that is used to group values together.

How data is represented?

The output data (输出数据) is a also set of <key, value> pairs.

e.g.: (key= CN, value = Shenzhen) (key= CN, value = Beijing)

A key (键) is some information that is used to group values together.

MapReduce

- MapReduce is a programming model (编程模型)
- It is inspired by the Map and the Reduce operations of the LISP programming language.
- It is designed to process **large datasets** on computing clusters (the **cloud** $\vec{\Xi}$).
- It is often used with the Java language.
- A programmer has to define map() and reduce() functions



A simple example

Consider that we want to count how many times each word appear in a very large text document.

"Hello world, bye world, Hello cloud, goodbye cloud"

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A simple example

The **master instance** first splits the data into **M** data blocks.





- All instances work in parallel.
- Consider the first instance. It reads its data.
- It creates a <key,value> pair (键值对) for each word that it reads. A key (键) is a word and the corresponding value (值) is the number 1.



"Hello world, bye world"

<Hello, I> <World, I> <Bye, I> <World, I>

- Some words like "World" appear multiple times in the result.
- All values that have the same key are grouped together.



- Consider the <u>second instance</u>.
- The second instance reads its data.
- It creates a <key,value> pair for each word that it reads. A key is a word and the corresponding value is the number I.



"Hello cloud, goodbye cloud" <Hello, !> <Cloud, !> <Goodbye, !> <Cloud, !>

Then, the **second instance** groups all values that have the same key together.







A simple example

So until now, we have:



<Hello, I> <World, **2**> <Bye, I>





Next, the **reduce phase** will combine the local results found by all instances. \rightarrow

A simple example - reduce

- The **master instance** will start **R** reducing instances for combining results of mapping instances.
- In this example, only one reducing instance is used (instance D)





A simple example - reduce



This is the final result!



Reducinginstance



<Hello, 2> <Cloud, 2> <World, 2> <Bye, 1> <Goodbye, 1>



A simple example

This is the code for this example:

```
map(String key, String value):
    //key: document name; value: document contents
    for each word w in value:
    EmitIntermediate (w, "1");
```

Combine local results

```
reduce (String key, Iterator values):
    // key: a word; values: a list of counts
    int result = 0;
    for each v in values:
    result += ParseInt (v);
    Emit (AsString (result));
```



(1) An application starts a master instance and M worker instances for the Map phase and, later, R worker instances for the Reduce phase.


(2) The master split (分配) the input data in M segments (parts).



(3) Each Map instance reads its input data segment and processes the data



(4) The local results are stored on the local disks of the computers where the Map instances are executed.





(6) The final results are written by the Reduce instances to a shared storage (共享存储)



(7) The master instance monitors the Reduce instances and. When all of them have finished, it is the **END**.



More details

The data is usually split in blocks of 16 MB to 64 MB (megabytes - 兆字节).

 The number of instances can be a few to hundreds, or thousands of instances.

What if some instances crashes? ightarrow



What happen if an instance fails?

- Fault-tolerance (容错): to ensure that a task is accomplished properly even if some machines stop working.
- The master instance asks each worker machine about their state (idle 空闲状态, in-progress 正在进行, or completed 完成任务) and identity.
- If the worker machine does not respond, the master instance considers that this machine's sub-task has failed.



What happen if an instance fails?

- A task in progress (正在进行) on a failed worker is set to idle (空闲状态).
- The task can then be given to another worker (computer).
- The master writes takes of note of the tasks that have been completed.
- The **data** is stored using the **GFS** (Google File System).

What is a typical MapReduce machine in a cluster?

According to the book, in 2012, a typical computer for experimenting with MapReduce has the following characteristics:

- dual-processor x86 running Linux,
- 2–4 GB of memory,
- Network card: 100–1,000 Mbps.
- Data is stored on IDE 7 disks attached directly to individual machines.
- The file system uses replication (复制)

What is a typical MapReduce machine in a cluster?

- A cluster consists of hundreds or thousands of machines.
- It provides availability (可利用性) and reliability (可靠) using unreliable hardware.
- The **input data** is stored on the **local disk** of each instance to reduce communication between computers.



A second example

Task: analyze a text to count how many words with I letters, with 2 letters, with 3 letters, with 4 letters...

"Hello world, bye world, Hello cloud, goodbye cloud"

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Α

All instances work in parallel.

Consider the <u>first instance</u>. The first instance reads its data. It creates a <key,value> pair for each word that it reads. A key is the number of letters in the word and the value is the word.





<5, Hello> <5, World> <3, Bye> <5, World>

Some words like "World" appear multiple times in the result.

All values that have the same key are grouped together.



Note: value having the same key are automatically grouped

Consider the <u>second instance</u>.

The second instance reads its data. It creates a <key,value> pair for each word that it reads, where a key is a number of letters and the corresponding value is a word.



"Hello cloud, goodbye cloud"

<5, Hello> <5, Cloud> <7,Goodbye> <5, Cloud>



Then, all values that have the same **key** are grouped together.



"Hello cloud, goodbye cloud"

Note: value having the same key are automatically grouped

<5, Hello> <5, Hello, Cloud, <5, Cloud> Cloud> <7,Goodbye> <7, Goodbye> <5, Cloud>



A second example

So until now, we have:





<3, Bye> <5, World, World, Hello>

<5, Hello, Cloud, Cloud> <7, Goodbye>

Now, the **reduce phase** will take place to combine the local results found by each instance

A second example - reduce

- The **master instance** starts **R** reducing instances for combining results of mapping instances.
- In this example, only one reducing instance is used (instance D)





A second example - reduce

The result is shown below. It means that there is one word containing three letters, six words containing five letters, and one word containing seven letters.





- Consider a social network (社会网络) like Wechat, QQ, LinkedIn where you can be friend with other people.
- If you are a LinkedIn user and you view the LinkedIn page of a friend, the page will indicates how many friends you have in common.
- Illustration \rightarrow



- Suppose that we have a social network with five users: A,B,C,D,E
- We assume that friendship (友谊) is a bidirectional relationship (双向关系).
- In other words, if you are a friend of someone, s/he is also your friend.

 Assume that this is the friendship graph:



Assume that data about friendship between users is stored in a text file as follows:

- A -> B C D
- B -> A C D E
- C -> A B D E
- D -> A B C E
- E -> B C D





The data file will be split and sent to various **mapping instances**.

Mapping instances will process each line that they receive as follows: \rightarrow



by combining A with each of his friend.

A third example - map The second line B -> A C D E is transformed as:

KeyValue $(A B) \rightarrow A C D E$ $(B C) \rightarrow A C D E$ $(B D) \rightarrow A C D E$ $(B E) \rightarrow A C D E$

A third example - map The third line C -> A B D E is transformed as:

 Key
 Value

 $(A \ C) \rightarrow A \ B \ D \ E$ $(B \ C) \rightarrow A \ B \ D \ E$
 $(B \ C) \rightarrow A \ B \ D \ E$ $(C \ D) \rightarrow A \ B \ D \ E$
 $(C \ E) \rightarrow A \ B \ D \ E$

A third example - map The fourth line D -> A B C E is transformed as:

 Key
 Value

 $(A D) \rightarrow A B C E$
 $(B D) \rightarrow A B C E$
 $(C D) \rightarrow A B C E$
 $(D E) \rightarrow A B C E$

A third example - map The fifth line E -> B C D is transformed as:

KeyValue $(B E) \rightarrow B C D$ $(C E) \rightarrow B C D$ $(D E) \rightarrow B C D$

A third example - map (sort)

The values are then grouped by their key:

- (A B) -> (A C D E) (B C D)
- (A C) -> (A B D E) (B C D)
- (A D) -> (A B C E) (B C D)
- (B C) -> (A B D E) (A C D E)
- (B D) -> (A B C E) (A C D E)
- (B E) -> (A C D E) (B C D)
- (C D) -> (A B C E) (A B D E)
- (C E) -> (A B D E) (B C D)
- (D E) -> (A B C E) (B C D)

Furthermore, they are sorted (as above)

A third example - reduction

This data is then split and sent to reducers

- (A B) -> (A C D E) (B C D)
- (A C) -> (A B D E) (B C D)
- (A D) -> (A B C E) (B C D)
- (B C) -> (A B D E) (A C D E)
- (B D) -> (A B C E) (A C D E)
- (B E) -> (A C D E) (B C D)
- (C D) -> (A B C E) (A B D E)
- (C E) -> (A B D E) (B C D)
- (D E) -> (A B C E) (B C D)



Each reducer will intersect the list of value on each line:

The first line: $(A B) \rightarrow (A C D E) (B C D)$ will thus become: $(A B) \rightarrow (C D)$



Each reducer will intersect the list of value on each line:

The second line: $(A \ C) \rightarrow (A \ B \ D \ E) (B \ C \ D)$ will thus become: $(A \ C) \rightarrow (B \ D)$

and so on....

A third example – final result

The final result is: Key Value (A B) -> (C D)(A C) -> (B D)(A D) -> (B C)(B C) -> (A D E) $(B D) \rightarrow (A C E)$ (B E) -> (C D) $(C D) \rightarrow (A B E)$ (C E) -> (B D) (D E) -> (B C)
A third example – final result

The final result is:

Key Value $(A B) \rightarrow (C D)$ (A C) -> (B D)(A D) -> (B C) $(B C) \rightarrow (A D E)$ (B D) -> (A C E)(B E) -> (C D) $(C D) \rightarrow (A B E)$ (C E) -> (B D)(D E) -> (B C)

Having calculated this information, we know the friends in common between any pairs of persons.

For example: A and D have the friends B and C in common

A third example - conclusion

 In this example, we have explained how the MapReduce framework can be used to calculate common friends in a social network.

• Why doing this?

- Big social networks such as LinkedIn have a lot of money.
- By precaculating (预先计算) information about common friends, a social network can provide the information more quickly to users.
- This can be recalculated every day.

• 4.8-CLOUD FOR SCIENCE AND ENGINEERING



4.8

Cloud for science/engineering

- In the last 2000 years, science was mostly empirical.
- In recent decades, computational science (计算科学) has emerged where computers are used to simulate complex phenomena.
- Science may now combine:
 - theory, experiment, and simulation (仿真)



Cloud for science/engineering

Generic problems involving data, in science:

- Collecting experimental data.
- Managing very large volumes of data.
- Building and executing models.
- Integrating data and literature.
- Documenting experiments.
- Sharing the data with others; data preservation for long periods of time.

Cloud for science/engineering

Example of large databases:

- The Chinese National Space
 Administration may collect huge amount of data about space using various equipment.
- The Chinese Meteorological Administration may collect huge amount of data about the weather.

The cloud is useful to analyze such large amount of data.

Biology research

- Cloud computing is very important for biology research.
 - **Computation of molecular dynamics** is CPU intensive.
 - Protein alignment (蛋白质序列) is data-intensive.

• An example \rightarrow

Biology research - example

An experiment carried out by a group from Microsoft Research illustrates the importance of cloud computing for biology research [223]. The authors carried out an "all-by-all" comparison to identify the interrelationship of the 10 million protein sequences (4.2 GB size) in the National Center for Biotechnology Information (NCBI) nonredundant protein database using *AzureBLAST*, a version of the *BLAST*²³ program running on the Azure platform [223].

Azure offers VMs with four levels of computing power, depending on the number of cores: small (1 core), medium (2 cores), large (8 cores), and extra large (>8 cores). The experiment used 8 core CPUs with 14 GB RAM and a 2 TB local disk. It was estimated that the computation would take six to seven CPU-years; thus, the experiment was allocated 3,700 weighted instances or 475 extra-large VMs from three data centers. Each data center hosted three *AzureBLAST* deployments, each with 62 extra-large instances. The 10 million sequences were divided into multiple segments, and each segment was submitted for execution by one *AzureBLAST* deployment. With this vast amount of resources allocated, it took 14 days to complete the computations, which produced 260 GB of compressed data spread across more than 400,000 output files.

Using 3,700 instances, a task that would took about 7 years on a single computer was done in 14 days!

ADDITIONAL INFORMATION

Introduction

- Last week, we talked about MapReduce.
- MapReduce is a model to create cloud applications.
- It is used for developing applications that can be used in the cloud.
- It is called MapReduce because there are two steps called "Map" and "Reduce".



Introduction

- MapReduce is a popular model.
- There are many other models for developing cloud applications.
- For example:
 - Apache Spark
 - Apache Storm

```
•••
```





- **Spark** is more complicated than **MapReduce**.
- Spark offers more than 100 operators to transform data.
- Spark can be used with the Java, Python and Scala programming languages (编程语言).





• A problem of MapReduce is that it reads and write data many times to the storage 存储 (before and after each Map or Reduce operation).



• This can make a cloud application **slower**.

• Solution:

- Using **Spark**, data can be kept in memory.
- In other words, data is not read and written many times.
- Spark can read and transform data. However, Spark is "lazy". It only read and transform data when an action needs to be performed on the data.





- When Spark transforms data, the data is then stored in a structure called:
 Resilient Distributed Dataset (RDD).
 Resilient = 能复原的
 Distributed = 分布式
 Dataset = 数据
- All the transformations that are applied to data are remembered so that a dataset can be recovered if some failure happen.







Conclusion

- In this part, I have presented the MapReduce model, which is widely used for cloud computing.
- The first assignment is announced today.



http://philippe-fournier-viger.com/COURSES/CLOUD/



References

- Chaptre 4. D. C. Marinescu. Cloud Computing Theory and Practice, Morgan Kaufmann, 2013.
- <u>http://stevekrenzel.com/finding-friends-with-</u> <u>mapreduce</u>
- <u>https://hadoop.apache.org/docs/r1.2.1/mapred_t</u> <u>utorial.html</u>